

Comparing schedule generation of VSIDS against CPRU for RCPSP-t solvers

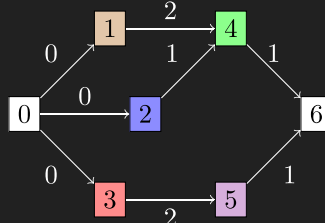
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1 Background

RCPSP is a class of problems where some **tasks**, which require **resources**, must be scheduled. RCPSP-t is a subproblem of RCPSP, where the resources produced and taken vary over time. The goal is to **minimize** the time from start to end [1]. CPRU is a heuristic designed for RCPSP-t and uses the **critical path** and **resource utilization**.

Precedence graph with durations:



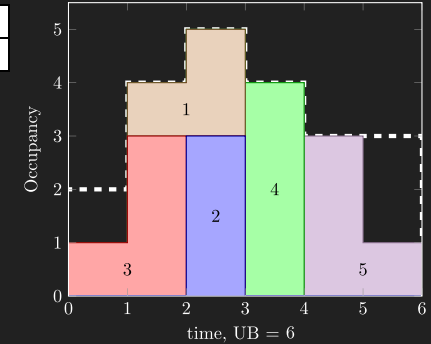
Resource capacities:

$B_{1,0}$	$B_{1,1}$	$B_{1,2}$	$B_{1,3}$	$B_{1,4}$	$B_{1,5}$
2	4	5	4	3	3

Activity durations and demands:

	0	1	2	3	4	5	6
p_i	0	2	1	2	1	2	0
$b_{i,1,0}$	-	1	3	1	4	3	-
$b_{i,1,1}$	-	2	-	3	-	1	-
$b_{i,1,2}$	-	-	-	-	-	-	-

An optimal solution:



2 Research Question

- How does selection randomization for CPRU compare in **runtime** and **average deviation from optimal** to a version without?
- Can CPRU be improved by using an **adaptive heuristic** calculation?
- How does CPRU compare against meta-heuristics like **VSIDS**[2]?

3 Methodology

RCPSP-t can be solved using **constraint programming (CP)**. In CP, variables which may take a range of values and constraints which enforce relations are used. A solver assigns each variable a value such that the constraints hold.

CPRU calculates the **critical path** by finding the longest path between the task and the end on the precedence graph. The **resource utilization** is the amount of resource a task and its successors take, relative to the time they could be scheduled in. CPRU only looks at a **subset** of the variables when selecting a variable.

4 Experiments & Conclusion

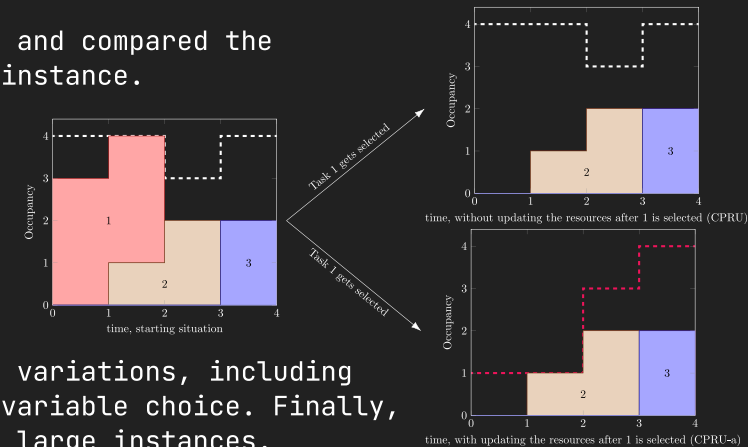
We evaluated 2 different variations of CPRU: CPRU without randomization and CPRU with the **resource utilization** updating as shown in the figure below (CPRU-a). We tested the variations and CPRU against VSIDS running them for **60 seconds**. We also tested them running for a **maximum of 10 schedules** within 9 hours.

We used the testing set by Hartmann [1] and compared the result against the lower bound on each instance.

Test results:

algorithm\test set	60-sec J30	60-sec J120	10-schedules J30
CPRU	44%	52%	43%
CPRU-nt	44%	52%	43%
CPRU-a	44%	52%	43%
VSIDS	44%	117%	92%

We concluded there was no significant performance difference between the CPRU variations, including showing CPRU-a had little influence on variable choice. Finally, we concluded CPRU outperforms VSIDS for large instances.



References

- [1]: S. Hartmann, Project scheduling with resource capacities and requests varying with time: a case study. Flexible Services and Manufacturing Journal, 25(1-2):74-93, 06 2013
 [2]: M.W. Moskewicz et al. Chaff: engineering an efficient sat solver. In Proceedings of the 38th Annual Design Automation Conference, DAC '01, pages 530-535, New York, NY, USA, 2001. Association for Computing Machinery.

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