

MODELLING CYCLIC STRUCTURES IN AGDA

COINDUCTIVE FORMALIZATIONS OF LINEAR TEMPORAL LOGIC

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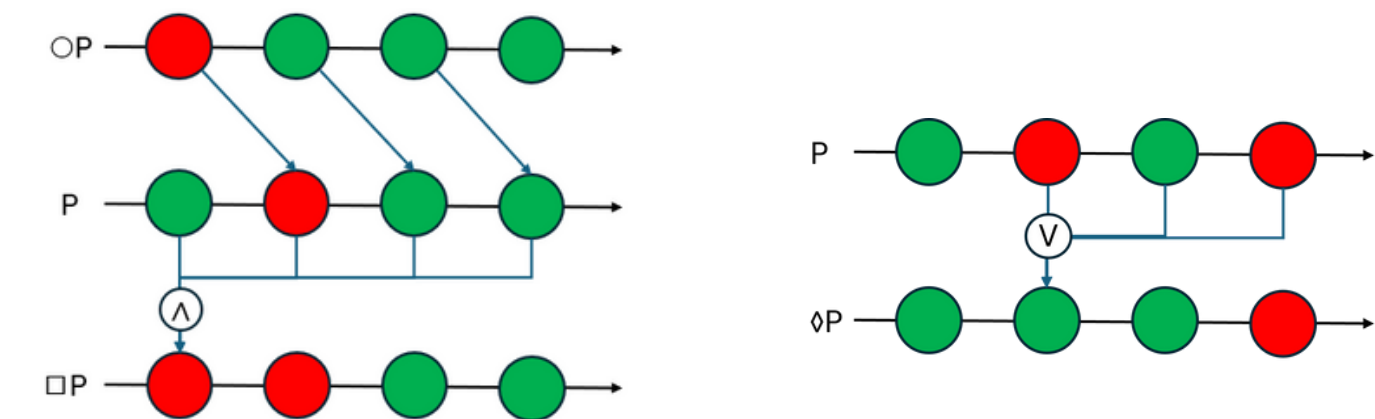
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1. BACKGROUND

LINEAR TEMPORAL LOGIC (LTL)

- Truth value of propositions varies in time.
- Pointwise lifting of classical operators and temporal modalities.
- A proposition is satisfied if it is currently true.
- A proposition is valid if it is true at any time point.
- *Kripke structure*: infinite sequence of valuations of propositional constants.



COINDUCTION

- *Corecursion*: Construct objects of a coinductive data type, by constructing the observations.
- Agda supports three types of coinduction:
 - Musical* coinduction: compose delays with $\#$ and force them with \flat operators.
 - Guarded* coinduction: use coinductive records and copattern matching to define coinductive types.
 - Sized types*: tracks depth of recursion structurally by embedding types with a size.

2. RESEARCH QUESTION

How can linear temporal logic be formalized in Agda?

- How can coinduction be used to model LTL formulas?
- What limitations of Agda arise in the process of modeling and reasoning about temporal logic?

3. ENCODINGS

DEEP EMBEDDING OF SEMANTICS

- Encode the syntax explicitly as an inductive type.
- Encode Kripke structures as streams of states.
- Encode the semantics using pattern matching.

```

_⊢_ : (KripkeStructure n) → LTLFormula {n} → Set
K ⊢ (Var x₁) = (head K) x₁ ≡ true
K ⊢ (Neg x₁) = (head K) x₁ ≡ false
K ⊢ (◯ P) = (tail K) ⊢ P
K ⊢ (□ P) = K ⊢ P
K ⊢ (◇ P) = Σ ℕ (eventually K P)

```

```

record ⊢_ (K : KripkeStructure n) (P : LTLFormula) : Set where
  coinductive
  field
    TrueNow : K ⊢ P
    TrueInFuture : (tail K) ⊢ P

```

Listing 1: Definition of satisfaction for the temporal operators, propositional constants, and negation.

4. EXPERIMENTS

1. **Soundness**: Derived an

axiomatization of LTL to prove correctness of the formalization.

2. **Absorption theorems**: Proved

identities exploring the interaction between \Diamond and \Box temporal modalities.

3. **Towers of Hanoi**: Encoded Towers of Hanoi as a state system and prove properties of the system.

$(A_2) K \models (A \rightarrow (B \rightarrow C)) \rightarrow (A \rightarrow B) \rightarrow A \rightarrow C$	$(Itl3) K \models \Box A \rightarrow A \wedge \Box \Box A$	1. $\Box \Box P \Leftrightarrow \Box P$
$(A_3) K \models (\neg B \rightarrow \neg A) \rightarrow A \rightarrow B$	$(MP) K \models A \rightarrow (A \rightarrow B) \rightarrow B$	2. $\Diamond \Diamond P \Leftrightarrow \Diamond P$
$(Fun) K \models \neg \Box A \Leftrightarrow \Box \neg A$	$(N\Box) K \models A \rightarrow \Box A$	3. $\Box \Diamond \Box P \Leftrightarrow \Diamond \Box P$
$(KO) K \models \Box (A \rightarrow B) \rightarrow (\Box A \rightarrow \Box B)$	$(Ind) K \models (A \rightarrow \Box A) \Rightarrow K \models (A \rightarrow \Box A)$	4. $\Diamond \Diamond \Diamond P \Leftrightarrow \Diamond \Box P$

Axiomatization of Linear Temporal Logic

Absorption Theorems

Encoding	Soundness	Absorption Theorems	Towers of Hanoi
Shallow embedding	✓	✓	✓
Deep embedding	✓	✓	✗

Table 1: Summary of the results of the experiments.

5. AGDA LIMITATIONS

- **Documentation**: Documentation for certain topics is missing or incomplete.
- **Error messages**: Oftentimes error messages are confusing and do not point to the actual problem.
- **Termination checking**: The termination checker fails to identify structurally decreasing parameters.

6. FUTURE WORK

- Explore **extensions** of the LTL we considered.
- Investigate formalizations of LTL using other types of coinduction.