

An analysis of Structured Encryption (StE) compared to other computation techniques on encrypted data

1. Motivation & Background

Outsourcing data and computation is very common nowadays with cloud environments. However, sensitive data with high privacy requirements, such as medical records, cannot be stored as-is on a remote server. It must be encrypted. Classical encryption techniques, such as AES, do not allow the server to make any kind of operation on the data, essentially making the computation power of the cloud useless.

Structured Encryption is a set of techniques that stores structured data encrypted on an untrusted server, while still allowing the client to outsource operations securely.

2. Research Question

What are the capabilities and limitations of modern Structured Encryption technologies?

How does Structured Encryption perform compared to Oblivious RAM (ORAM), Fully Homomorphic Encryption (FHE), Trusted Execution Environments (TEEs) and Secure Multiparty Computation (SMC)?

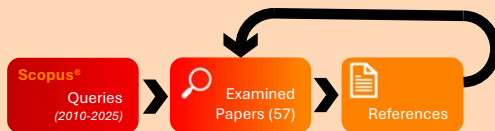
State of the art technologies

Threat Models / Limitations

Characteristics Comparison

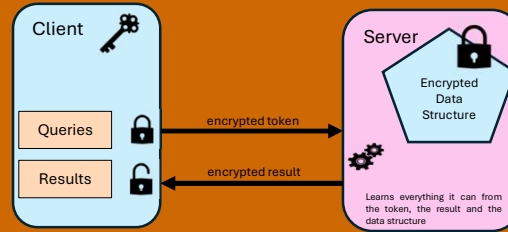
3. Methodology

Snowball Sampling through Backward Reference Searching



4. Technologies & Characteristics

- SSE (Symmetric Searchable Encryption)**: symmetric-key cryptographic technique allowing a user to search over encrypted labeled data stored on an untrusted server, without revealing the data or the search terms. [1]
- StE**: generalization of SSE to any data structure. [2]
- Semi-honest model (honest-but-curious)**: the server performs the operations honestly but learns as much info as possible in the process

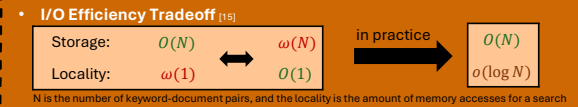


4.1 Functionality and Use Cases

- Labelings**: boolean queries [6] (AND, OR and NOT), add / delete operations [5], substring & wildcard queries
 - Numerical data**: range queries [7]
 - Graphs**: adjacency, shortest path and focused subgraph queries [2]
 - Concept and Knowledge Graphs (for AI models) [8]
- ⇒ Usages: MongoDB's Queryable Encryption, Cipherstash, Crypteron

Encrypted Databases (SQL queries) [3]

4.2 Efficiency



- Parallelism**: improvement with sufficient cores: $O(r) \rightarrow O(\log N)$ [4]
 r is the size of the query result
- Concurrency**: lock-free concurrent update operations

4.3 Security

- CKA1 (Chosen Keyword Attack)**: non-adaptive security [9]
- CKA2**: adaptive security, attacker can choose keyword after seeing the encryption. This is the standard for StE schemes.
- UC (Universally Composable)**: a cryptographic protocol remains secure even when composed with other arbitrary protocols. Very strong security but often too heavyweight.
- Leakage**: information leaked to the server in the setup or query process
 - Volume Leakage**: the size of the result of a query
 - Access pattern leakage**: which data corresponds to the query

Attacks

- File-injection attacks on dynamic multimaps: 1% of file injected provides 40% efficiency (2016)
- Inference attacks on keywords: exponential in the amount of keyword and requires 70% of dataset knowledge to become slightly successful. (2012 & 2015) [10]
- Inference attacks on range queries: only $O(N \log N)$ queries needed to reconstruct the entire data (2017 & 2018)
- Encrypted databases attacks on column-equality leakage (2024)

Countermeasures

- backward and forward private schemes (addition/deletion do not reveal information about previous and future queries)
- access-pattern hiding schemes (ORAM-based schemes, from 2016) e.g. TWORAM, EMMavl
- volume-hiding schemes (from 2018) e.g. d-DSE, XorMM [16] (polylog efficiency overhead)

ORAM = technique that hides memory access patterns. [12]

5. Comparison MPC = method that allow multiple parties to jointly compute a function over their inputs without revealing them. [13]

Similarities with StE	Differences with StE
Threat model: Semi-honest	Generic data access instead of structured data
Sublinear (logarithmic) efficiency	The client must keep track of the data structure
Client-Server architecture	Does not leak the access pattern in queries

- ORAM can be used as a component in StE schemes to hide the access pattern and provide protection against inference attacks, by adding a logarithmic overhead to the query complexity. e.g. TWORAM

FHE = encryption scheme allowing computation to be directly performed on encrypted data. [11]

Similarities with StE	Differences with StE
Client-Server architecture	Can perform any computation
Non-interactive scheme	Does not leak any information
	Polynomial efficiency overhead

- FHE can be used to perform the same task as StE schemes with no leakage, but with impractical efficiency

Similarities with StE	Differences with StE
Threat Model: Semi-honest (sometimes malicious)	Can perform any computation
Many different protocols for many use cases	Does not leak any information
Both can perform Private Set Intersection (PSI)	Client-Server and Distributed architecture

- StE provides the most efficient solution for Updatable PSI, which is a typical MPC problem. MPC remains the best solution for static PSI.

TEEs = secure area within a processor ensuring that data and code are protected. [14]

Similarities with StE	Differences with StE
Client-Server architecture (cloud environment)	Can perform any computation
Practical efficiency (near native)	Threat Model: malicious server software
	Hardware-based

6. Conclusion

- StE schemes provide practical efficiency and functionality and is used in real-world DBMS like MongoDB
- StE leaks information to the server, and schemes not providing the most recent security features like forward/backward privacy, volume-hiding and access pattern hiding, can be attacked easily.
- Future research should focus on new attacks and study of the complex leakage profile of StE schemes.

- In contrast with the other techniques, StE schemes do leak some information, to achieve practical efficiency.
- Future research could focus on the creation of a precise benchmark to compare these techniques on practical scenarios.

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