

Code Extraction From a Dependently Typed Language - “How Effectively Does Forth Serve as a Target Language for Extraction From Agda?”

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1. What Is Agda?

- Pure functional language
- Dependently typed
 - A dependent type are type whose definition depends on another value.
- E.g. $(Vec\ 5)$ could represent an array of length 5

```
data Vec (A : Set) : Nat → Set where
  []      : Vec A zero
  _::__   : {n : Nat} → A →
           Vec A n → Vec A (suc n)
```

- Used to provide guarantees about code and write proofs
- Lazy evaluation
- Some code is unused after compilation

2. What Is Forth?

- Stack based, imperative language
- Untyped
- Values are pushed onto the stack instead of being passed as arguments
- Words (functions) pop values off the stack to use them

```
5 6 ok
.s <2> 5 6 ok
+ ok
.s <1> 11 ok
```

- Strict evaluation

3. The Problem

- How to translate Agda to Forth?
- How will it perform compared to other compilers? (Haskell, Scheme)
- Is it even possible to implement every feature of Agda in Forth
 - If not, how much can be?

4. Results

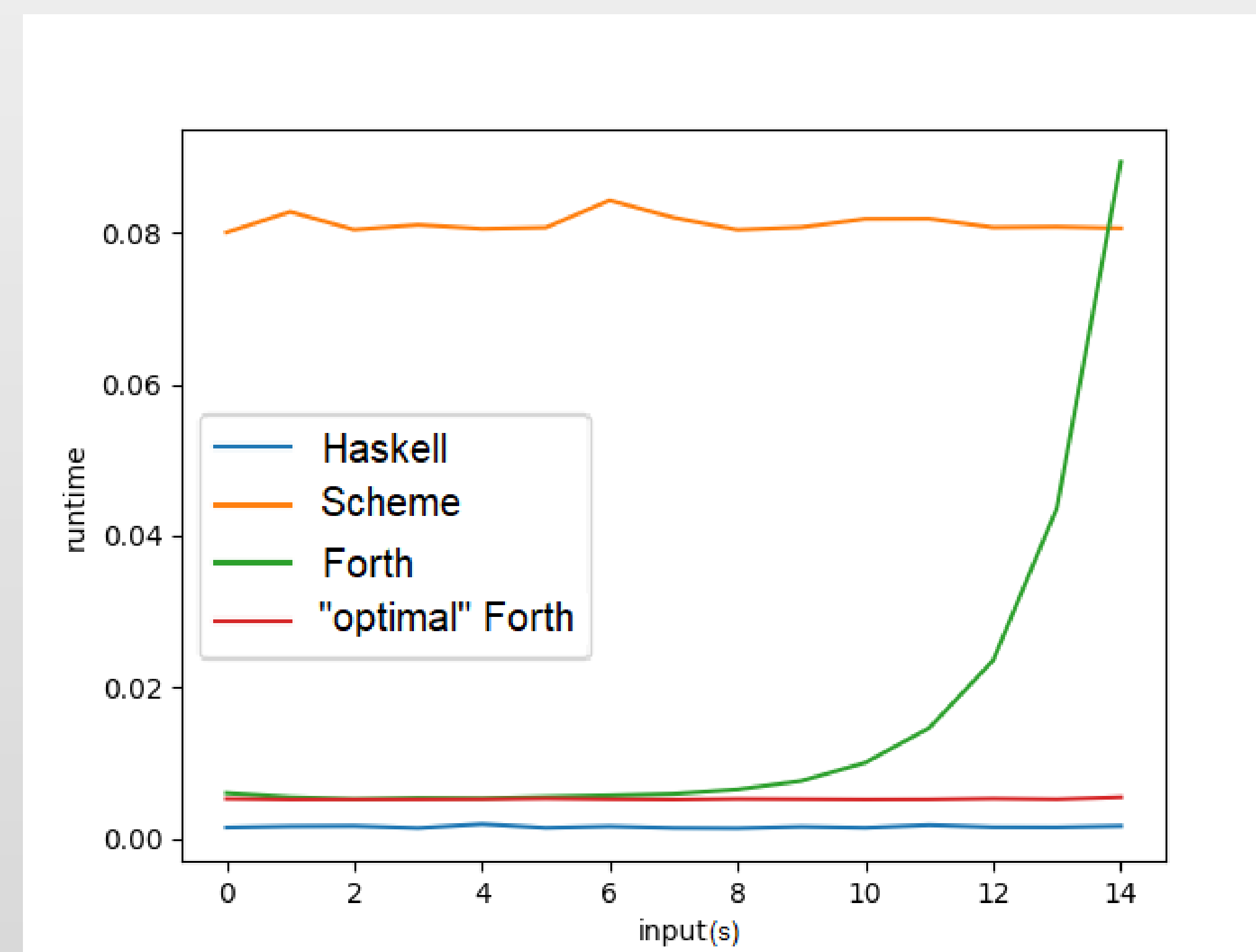


Fig. 1: Runtimes for consume function with inputs 0 - 15

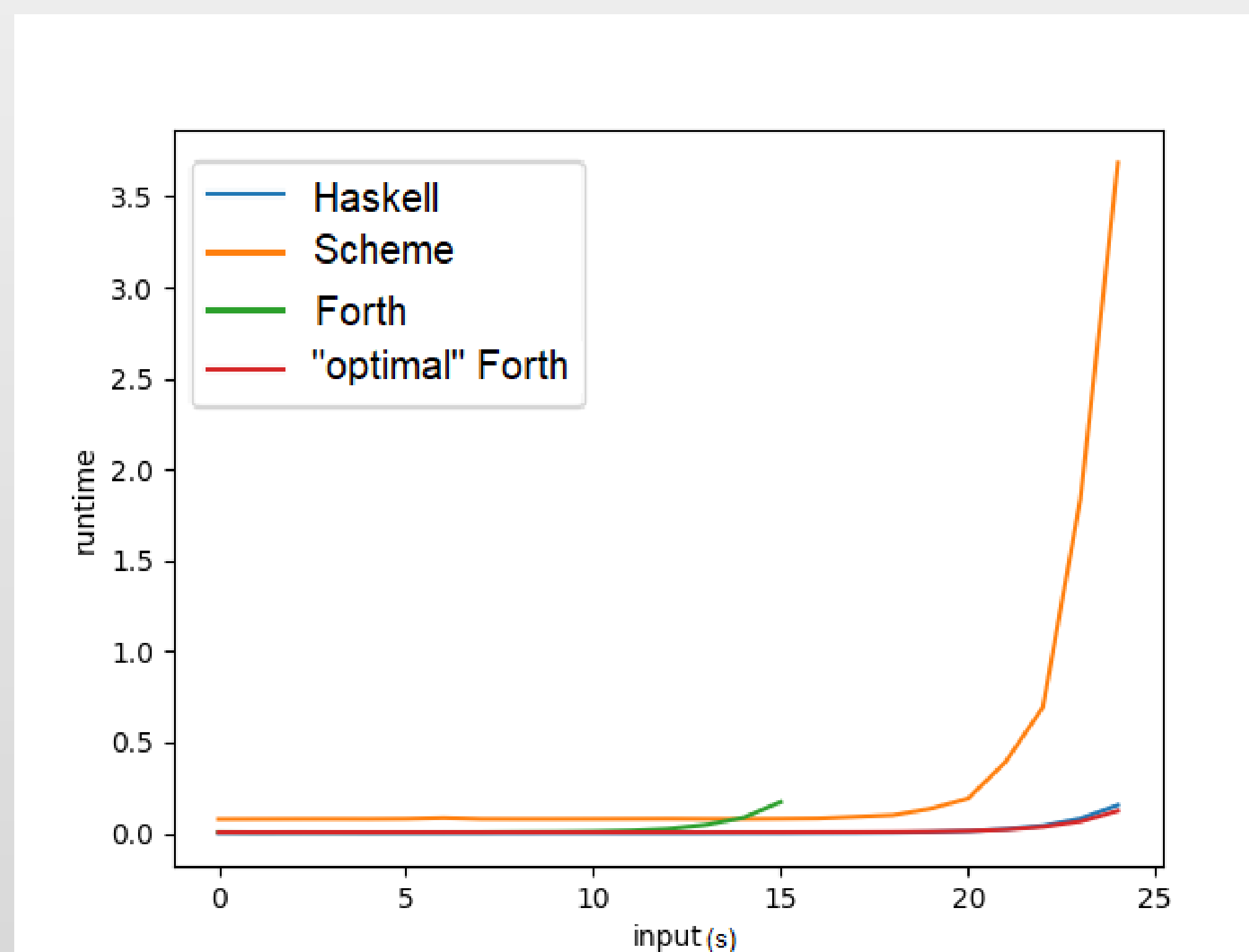


Fig. 2: Runtimes for consume function with inputs 0 - 24

5. Conclusions

- Forth is not a very good target language
- Lack support for too many features of Agda
- Improvements to my compiler could greatly improve performance
 - Improve thunk handling
 - Consider different Forth implementations

[1] Agda GHC Backend:
<https://agda.readthedocs.io/en/v2.6.2.1/tools/compilers.html#ghc-backend>

6. Example Translations

```
consume : Nat → Nat
consume zero = zero
consume (suc n) = consume n
```

Fig. 3: consume function in Agda

```
defer consume
...
ENDFLAG :noname { a }
a 0 0 pointer ! obj=
if
  wildcards 0 pushArrayN { }
  0
else
  1 a sub pass pass { b } b consume pass
then
;
create XTconsume fillHere
:noname XTconsume foldThunks ; is consume
```

Fig. 4: consume function in Forth

```
ite : {A : Set} → Bool → A → A → A
ite true x y = x
ite false x y = y
```

Fig. 5: ite (if, then else) function in Agda

```
defer ite
...
ENDFLAG :noname { d c b a }
b if c else d then ;
:noname { c b a e }
c b a e pass pass pass ;
:noname { b a f }
b a f pass pass ;
:noname { a g }
a g pass ;
create XTite fillHere
:noname XTite foldThunks ; is ite
```

Fig. 6: ite (if, then else) function in Forth