

1. Agda

- Dependently typed functional programming language.
- Can be used as a proof assistant.
- Formal verification.

2. Motivation

- How to ensure a refactoring is correct?
- Extensive testing increases confidence, but provides no explicit guarantees as opposed to a formal verification [1].
- Less attention for functional languages as opposed to class-based, object-oriented languages [2].

3. Refactoring tuples to records

- Tuples are less explicit compared to records.
- Record field accessors provide layer of abstraction.
- Not possible to distinguish tuples during refactoring when they have the same signature, but different representation.
- Is proving the correctness of tuple to record refactoring feasible?

```

-- Before: less expressive and extensible
lastName :: (String, String, Int) -> String
lastName (_, s, _) = s

-- After: more expressive and extensible
data Person = Person { initials :: String
                      , lastName :: String
                      , age :: Int
                      }
    
```

Listing 1. Tuple to record refactoring example based on work by Miran [3].

4. The Haskell-like language (HLL)

- Intrinsically-well-typed language constructed in Agda that forms the foundation of the refactoring operation and the proofs.

```

data _,_!_ : Ctx -> DataCtx -> Type -> Set where
...
tuple : TypeResolver Γ Γd ts
-----
→ Γ , Γd ⊢ tupleT ts
tLookup : Γ , Γd ⊢ tupleT ts → t ∈ ts
-----
→ Γ , Γd ⊢ t
recInst : (recDecl ts) ∈ Γd → TypeResolver Γ Γd ts
-----
→ Γ , Γd ⊢ recT ts
rLookup : Γ , Γd ⊢ recT ts → t ∈ ts
-----
→ Γ , Γd ⊢ t
    
```

Listing 3. Tuple and record-related constructors for the HLL.

```

data _,_!_ : Env Γ → (Γd : DataCtx) → (Γ , Γd ⊢ t) → Value t → Set where
...
!tuple : ReductionResolver γ Γd tr vs
-----
→ γ , Γd ⊢ tuple tr ⊔ tuple vs
!tLookup : γ , Γd ⊢ e ⊔ tuple vs
-----
→ γ , Γd ⊢ (tLookup e x) ⊔ poly-list-lookup vs x
!recInst : ReductionResolver γ Γd tr vs
-----
→ γ , Γd ⊢ recInst x tr ⊔ rec vs
!rLookup : γ , Γd ⊢ e ⊔ rec vs
-----
→ γ , Γd ⊢ (rLookup e x) ⊔ poly-list-lookup vs x
    
```

Listing 2. Big-step semantics for tuple and record-related constructs.

5. Apply refactoring to the HLL

- **Refactoring operation:** for all tuples, generate a record declaration and replace the tuple by a record instance of that declaration. All declarations are globally known.
- Make a sub-expression aware of its context by providing a trace. This helps to update existing record declaration lookups and generate new ones for situations similar to figure 1.

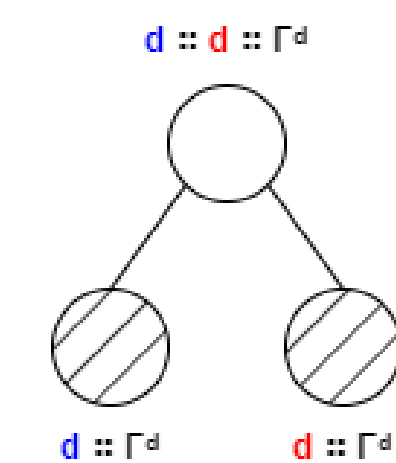


Figure 1. Example of the refactored declaration context at different levels of a language construct. Dashed circles indicate refactored tuples.

6. Proving correctness and conclusions

- Due to the use of an intrinsically-typed language, the refactoring operation is also a well-typedness proof.
- Prove that the refactoring replaces **all** tuples by mapping a refactored expression to a construct that does not support tuples.
- Relation construct is used to show how the refactoring alters the evaluated value of an expression.
- Successfully constructing these proofs leads to a positive feasibility indication.
- Techniques can be reused for other functional programming languages that share the notion of tuples and records (e.g., Erlang).

7. References

- [1] D. Horpácsi, J. Kőszegi, and S. Thompson, "Towards Trustworthy Refactoring in Erlang," Electronic Proceedings in Theoretical Computer Science, vol. 216, pp. 83–103, 7 2016.
 [2] E. A. AlOmar, M. W. Mkaouer, C. Newman, and A. Ouni, "On preserving the behavior in software refactoring: A systematic mapping study," Information and Software Technology, vol. 140, p. 106675, 2021.
 [3] L. Miran, Learn you a Haskell for great good!: A beginner's guide, ch. Making Our Own Types and Type Classes. No Starch Press, 2012.