Modelling cyclic structures in Agda

1. Motivation

- Graph theory plays an important role in a multitude of different fields such as, but not limited to the field of **chemistry**, **sociology**, biology, operations research and even war science and is widely used in computer science as well.
- Various graph representations exist: a **popular** approach is through an **inductive** way.
- **Cycles** occur in graphs \rightarrow Inductive approach becomes less intuitive.
- **Solution**: the dual of induction, coinduction, which is not as well researched.

2. Background

Properties of graphs:

- isPresent: when provided a node and a graph, it checks if the node is present in the graph.
- hasPath: when provided a start node, an end node and a graph, it checks if the end node can be reached from the start node in the graph.
- hasCycle: when provided a graph, it checks if the graph contains a cycle.

Agda:

- Functional language
- **Total language**: all functions must terminate.
- **Dependently typed**: types can be indexed by objects of other types.
- \rightarrow Agda can be used as a **proof assistant**.
- **Coinduction**: Agda has three flavours of coinduction, guarded coinduction, musical coinduction and sized types.

3. Research question

How can graphs be modelled coinductively in Agda such that they are suitable for the proof assistant?

Sub-questions

- 1.What are different encodings of graphs in Agda and which are suitable for the proof assistant?
- 2.What properties of graphs can be proven using the various encodings (e.g. has node)?
- 3. What improvements should be made to Agda in order for it to handle modelling with coinduction more easily?

