

## 1. Motivation

- Baseline model assumes agents must share generated resources immediately.
- Real systems often allow delay:** batteries store energy, caches hold data, and queues postpone computation.
- Storage changes the decision problem: **agents now choose not only who to share with, but also when to share.**

**How does local storage affect strategic decision-making, equilibrium, and efficiency in a decentralized money-free exchange economy?**

## 2. Storage-Enabled Equilibrium

- In a traditional competitive market, equilibrium is usually described by three conditions: market clearance, budget feasibility, and utility maximization.
- In reciprocal exchange, agents do not use money. They trade directly using their own generated resources. Therefore, budget feasibility and market clearance collapse into one market clearance condition

**Storage-enabled feasibility**  $D_i(t) + S_i(t-1) = \sum_{j \in \mathcal{N}_i} x_{ij}(t) + S_i(t)$

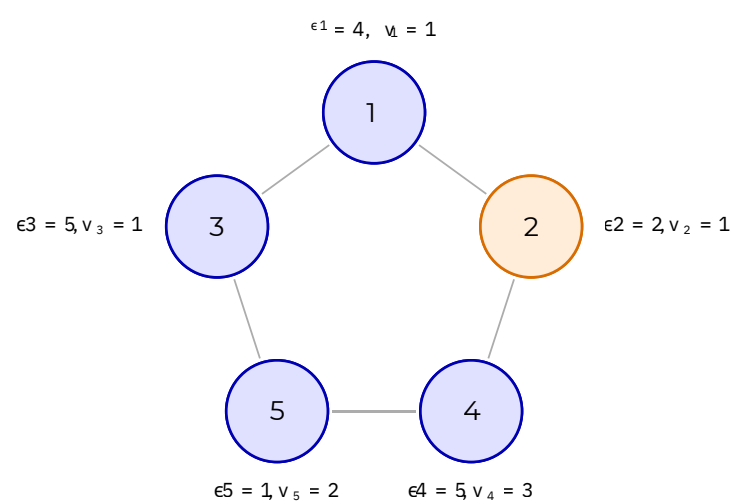
Each unit of available resource is either shared with a neighbour or carried forward in local storage.

**Intertemporal market clearance**  $\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \sum_{j \in \mathcal{N}_i} x_{ij}(t) = c_i$

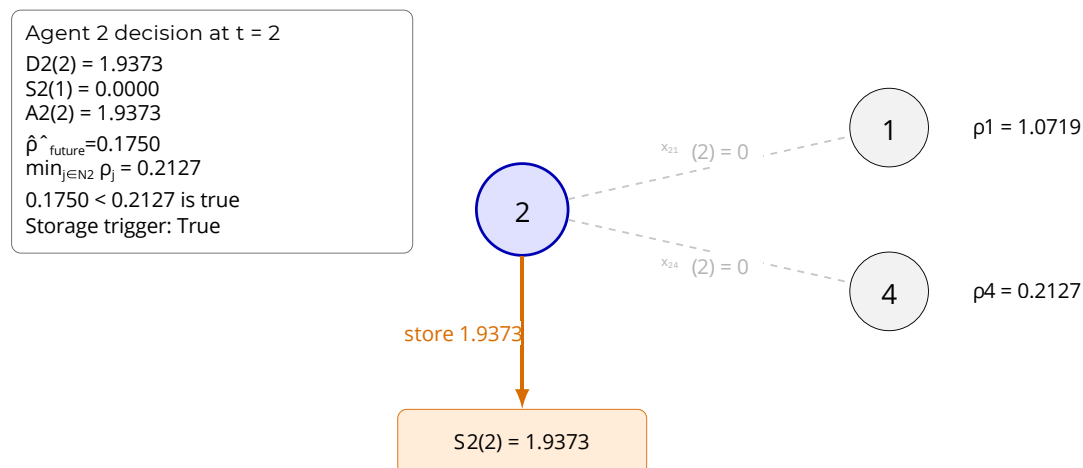
Because storage is bounded and lossless, it can delay exchange, but it cannot permanently remove resources from the market.

- The utility-maximization benchmark remains the same as in the baseline model. Storage changes when resources are exchanged, not how utility is defined.

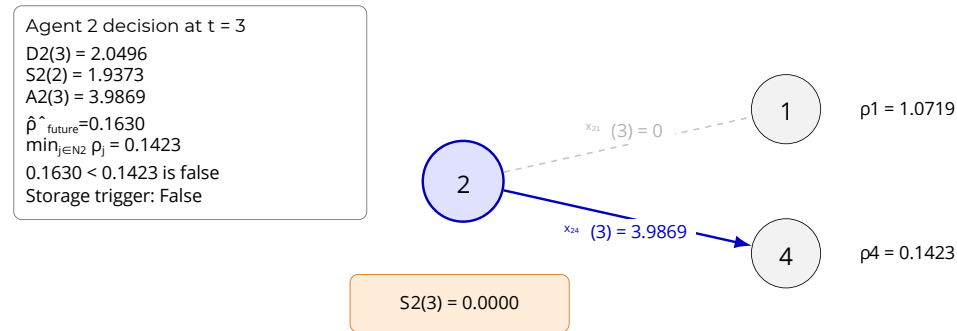
## 3. Greedy



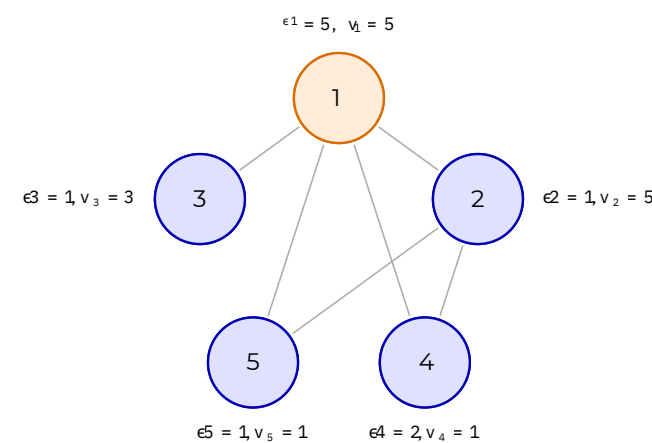
Pure greedy decision:  $t = 1 \rightarrow t = 2$



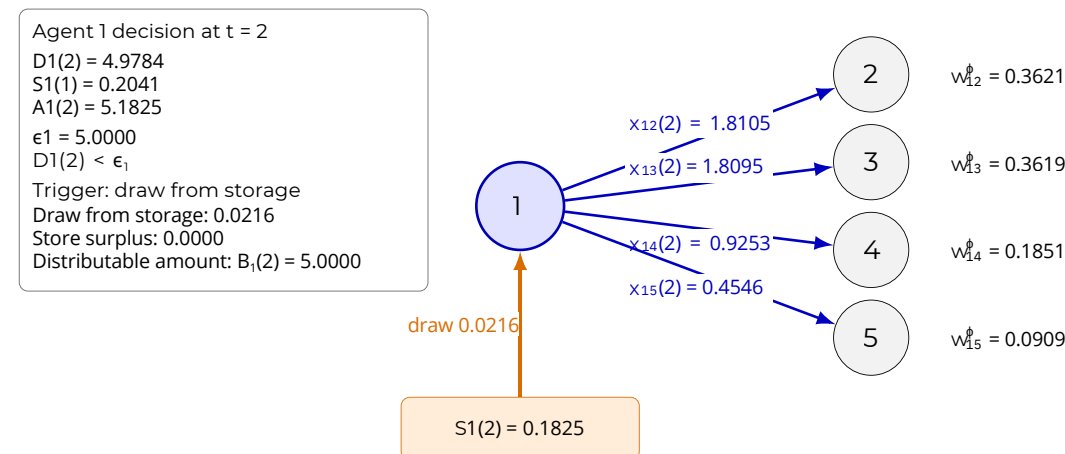
Pure greedy decision:  $t = 2 \rightarrow t = 3$



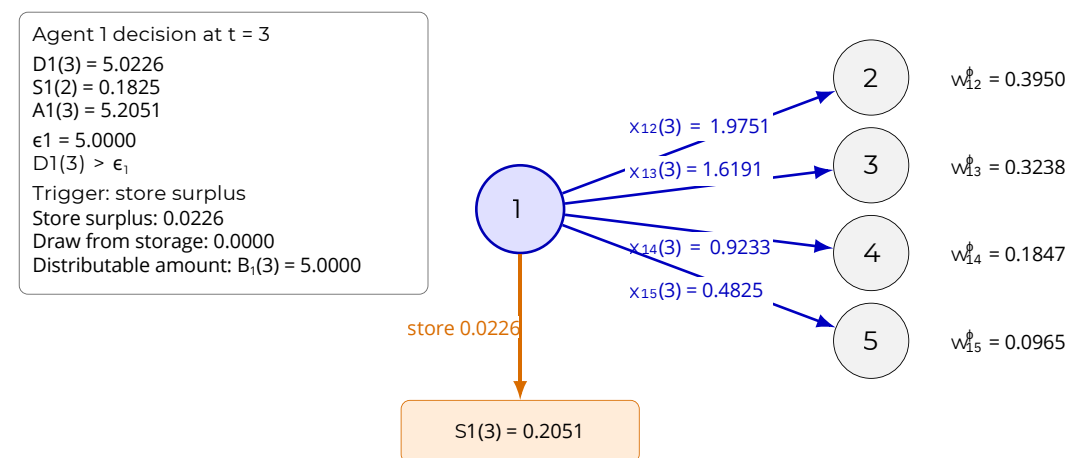
## 4. Proportional



Pure proportional decision:  $t = 1 \rightarrow t = 2$



Pure proportional decision:  $t = 2 \rightarrow t = 3$



## 5. Simulation Setup

Storage capacity	$c = 0, 0.5, 1, 2$
Strategies	$\phi$ proportional — $\pi$ greedy — $\psi_{hom}$ , $\psi_{het}$ mixed
Networks	complete — grid — random — scale-free — small-world
Measured outcomes	distance to benchmark — fairness — storage use

**Simulation Metadata**

$n = 20$  agents

$T = 400$  timesteps

$\sigma = 0.1$  generation noise

synchronous updates

$R = 5$  independent runs

**Storage capacity**

$S_{max} = c \cdot \frac{1}{|N|} \sum_{i \in N} c_i$

$c = 0$  is the zero-storage baseline.

- Main experiment: identify when storage improves the path toward equilibrium, and when it creates harmful delay.
- Greedy memory: how greedy agents behave when they trust recent market signals more than older ones.
- Higher volatility: check whether storage becomes more useful when resource generation is more volatile, i.e. when the standard deviation is higher.

## 6. Discussion

- Storage capacity is not always an improvement**
  - More storage gives agents more freedom to delay exchange, but delay is only useful if it leads to better future allocations.
  - In many cases, extra capacity just changes the timing without improving the final market state.
- Greedy storage can hurt, but not in every market instance.**
  - Greedy agents store when the future looks better than the current best neighbour.
  - This can work well in specific situations where waiting really does create a better exchange opportunity.
  - However, on average, the same behaviour mostly hurts.
- Recency changes storage usage, not outcomes.**
  - When greedy agents trust recent observations more, they store less and share sooner, but final loss and fairness do not change much.
- Volatility supports the smoothing-buffer interpretation**
  - When generation becomes noisier, proportional storage becomes more useful because it smooths short-term fluctuations around the long-term endowment.

## 7. Conclusion

- How can a decentralized sharing problem with local storage be formally modeled?

$$D_i(t) + S_i(t-1) = \sum_{j \in \mathcal{N}_i} x_{ij}(t) + S_i(t)$$

- What equilibrium concept is appropriate for a storage-enabled exchange economy?

$$\lim_{T \rightarrow \infty} \frac{1}{T} \sum_{t=1}^T \sum_{j \in \mathcal{N}_i} x_{ij}(t) = c_i. \quad \text{The utility-maximization benchmark remains the same as in the baseline model.}$$

- What is the strategy space of an agent that can choose between sharing resources immediately and storing them for future use?

$\phi$  proportional —  $\pi$  greedy —  $\psi_{hom}$ ,  $\psi_{het}$  mixed

- How can the model be simulated to evaluate the effect of storage capacity on efficiency, convergence, and fairness?

$$S_{max} = c \cdot \frac{1}{|N|} \sum_{i \in N} c_i$$

- How does storage change the path by which decentralized exchange approaches equilibrium?

It depends on the interaction between strategy type and network topology

## 8. Limitations and Future Work

- The experiments only cover part of the possible storage-enabled market space: fixed market size ( $n=20$ ), selected strategies, and selected network topologies.
- Storage is idealized as bounded and lossless. Real batteries, caches, or delayed jobs may lose value over time.
- Future work should test larger and more diverse markets and introduce storage decay or degradation to study the trade-off between sharing now and waiting.