## Maybe a List would be better?: Correct by construction Maybe to List refactorings in a Haskell-like language

#### CSE3000 Research Project

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### 1. Introduction

- Functional languages have referential transparency, i.e. **expressions** correspond to a single value
- Dependently typed languages, a form of functional programming, allow you to write and verify proofs using their type checker.
- You can have your cake and eat it too! Write your function and formally verify it.
- Goal of this RP is to apply this to refactorings in a **Haskell Like Language** (HLL) using the dependently typed language **Agda** [1]



safe-divM : Maybe Int -> Maybe Int -> Maybe Int safe-divM \_ (Just 0) = Nothing safe-divM (Just x) (Just y) = Just (x / y) safe-divM \_ \_ = Nothing safe-divL : [Int] -> [Int] -> [Int safe-divL \_ [0] = [] safe-divL [x] [y] = [x / y] safe-divL \_ \_ = []

Figure 1. Example refactoring

## 4. Primary takeaways

# Techniques generalize well to other (data-oriented) refactorings

- Extend\_Under\_ data type
- $\,\circ\,$  Keeps track of how structure of context changes
- $\mapsto_r$  relation
- Defines "valid" updates for data-oriented refactorings
- $\,\circ\,$  Definition in the case of closures weak but sufficient
- Intrinsic typing
- $\,\circ\,$  Implicitly handles well-typedness

## 2. Objectives

- 1. "Define a Haskell-like language in Agda"
- 2. "Write the refactorings of Maybe to List and formally verify its correctness"

## 3. Current work

#### Language Design choices

- Intrinsically-typed language
- $\,\circ\,$  Define type system and syntax at the same time
- De Bruijn indices
- $\,\circ\,$  Variable references are all unique
- Big step semantics

5. References

 $\circ\,$  Define pre-conditions and abstract small steps

case x of Nothing -> e\_n Just i -> e\_j

Figure 2. Problematic refactoring

#### Refactoring

- Intrinsic typing
- Refactoring function is proof of well-typedness
- Refactoring introduces new additions to Environment/Context
- Use Extend\_Under data type to keep track of changes

#### Proof

- Values can change post refactoring
- $\circ$  Define  $\mapsto_r$  relation to define how values should change
- Closure values may have different environment lengths and have different bodies
- Use "weak extensional equivalence" to define valid closure values post refactoring.
- $\circ\,$  "Equivalent" inputs give "equivalent" outputs

[1] Ana Bove, Peter Dybjer, and Ulf Norell. "A Brief Overview of Agda-A Functional Language with Dependent Types." In: TPHOLs. Vol. 5674. Springer. 2009, pp. 73–78.

[2] Philip Wadler, Wen Kokke, and Jeremy G. Siek. Programming Language Foundations in Agda. Aug. 2022. url: https://plfa.inf.ed.ac.uk/22.08/.