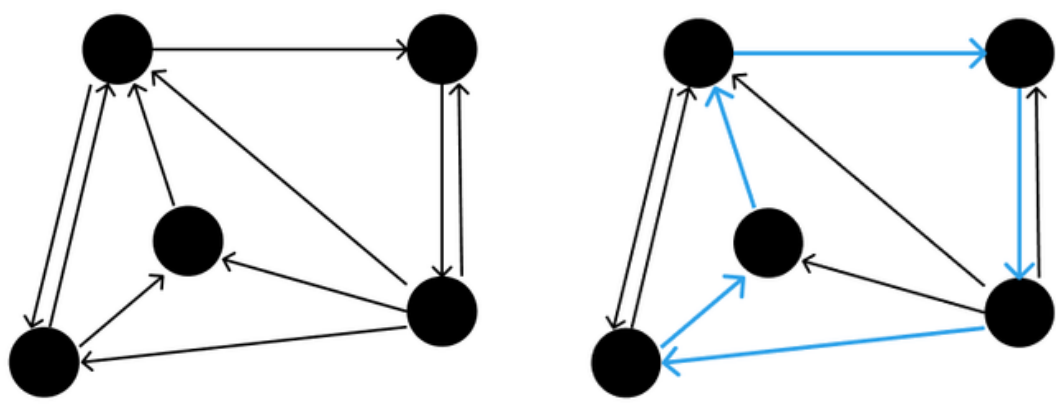


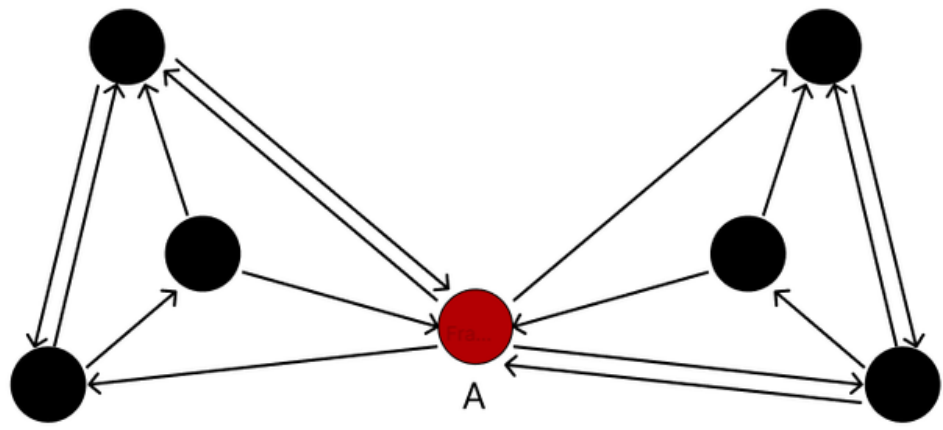
1. Motivation

Constraint Programming: a paradigm for solving combinatorial problems using constraints and propagation



Circuit Constraint: Find one Hamiltonian cycle through all vertices

- Existing propagators
- Prevent premature cycles
 - Local reasoning

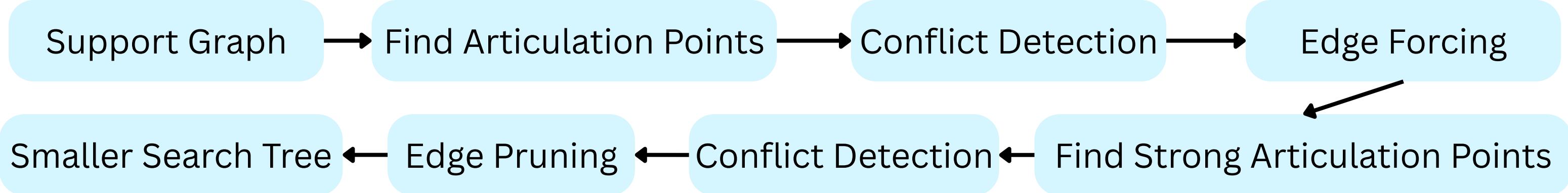


(Strong) Articulation points can detect global infeasibility early

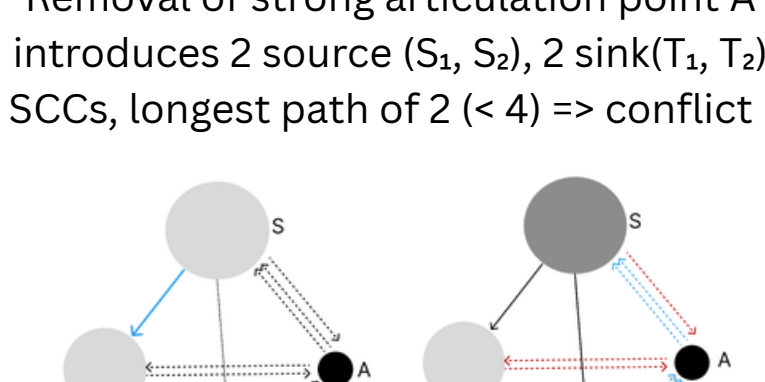
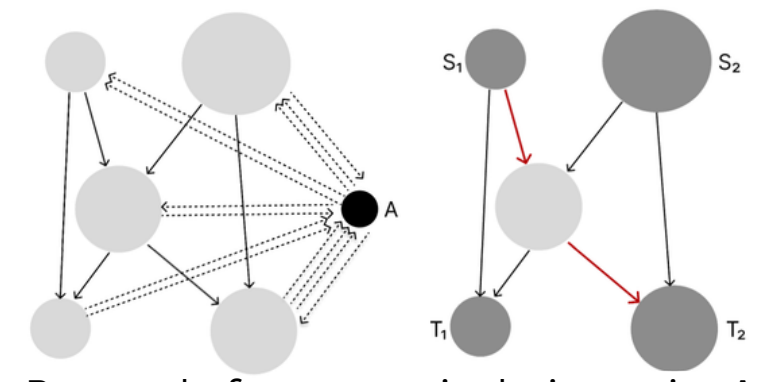
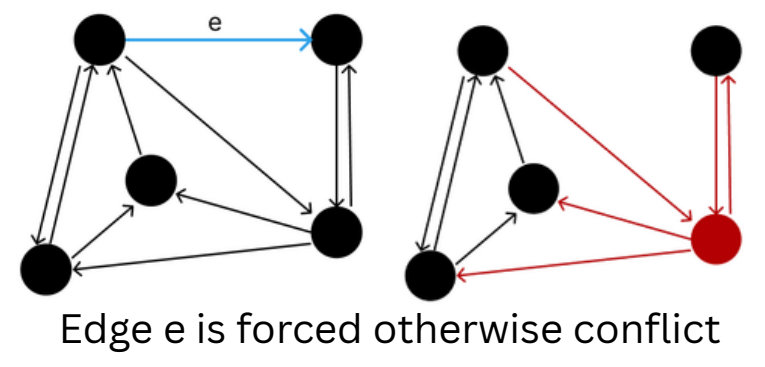
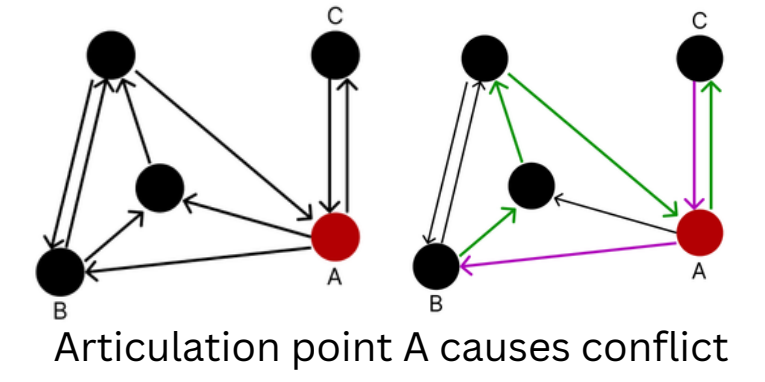
2. Main Idea

- Extend existing propagator
- Detect (strong) articulation points
 - Force edges
 - Prune edges
 - Detect conflicts earlier
- LCG Explanations
- Record all removed edges
 - Re-verify connectivity property

3. Propagation Workflow



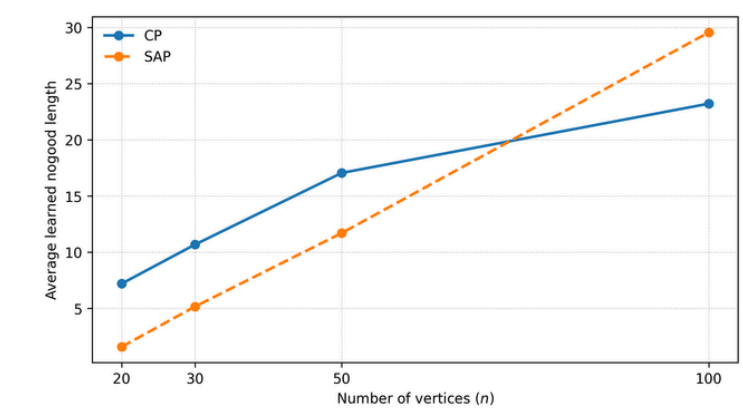
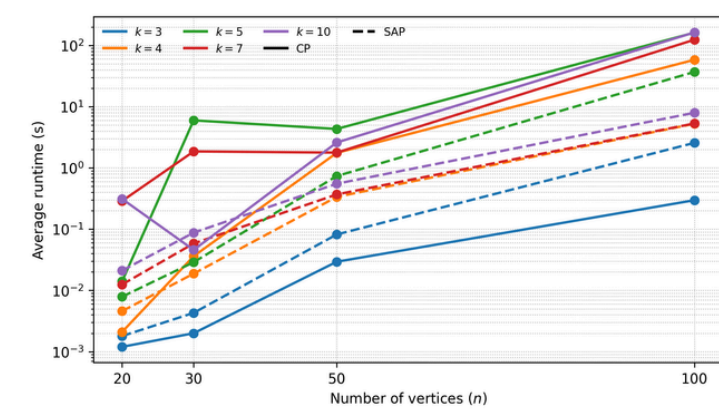
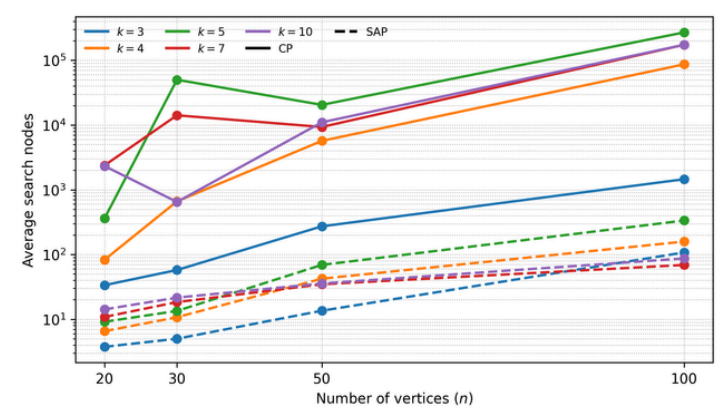
4. Reasoning



5. Experiments

- Pumpkin LCG solver
- $n = 20-100$ vertices
- varying k for different densities
- 400 benchmark instances
- Compared to cycle prevention

- Up to:
- **2510×** fewer search nodes
 - **7420×** fewer failures
 - **23×** faster runtime
 - **77%** fewer timeouts



- Search Nodes**
- Orders-of-magnitude search reduction
 - Benefits increase with graph size
 - Up to 2500× fewer nodes

- Runtime**
- Less search outweighs analysis overhead
 - Faster on most benchmark classes
 - Up to 23× speedup

- Learned Nogoods**
- Slightly larger explanations on large graphs
 - Expected due to global graph reasoning

6. Conclusion

- Connectivity bottlenecks provide powerful propagation
- SAP reasoning substantially reduces search effort
- Up to 2500× fewer search nodes
- Up to 7400× fewer failures

7. Future Work

- Incremental graph analysis
- Explanation minimization
- Adaptive activation based on graph density
- Evaluation on real-world routing benchmarks